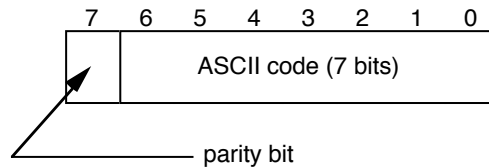


Review of ASCII character representation:



ASCII uses 8 bits to represent characters. Actually, only 7 bits are used to uniquely define the character and the 8-th bit (called the parity bit) is used for error detection. When used, the value of the parity bit depends upon the numbers of 1's in bits 0-7. For odd parity, bit 8 is set to make the total number of 1's in the byte an odd number such as 1 or 7. For even parity, bit 8 is set to make the total number of 1's in the byte an even number such as 0, 2 or 8.

Some useful ASCII character codes:

character	ASCII code (in hex)
/	2F
0	30
1	31
2	32
8	38
9	39
:	3A
;	3B
@	40
A	41
B	42
Z	5A
[5B
\	60
a	61
z	7A
{	7B
etc.	

EXAMPLE: PARITY PROGRAM

The correct way to design a program is by starting with your inputs, outputs and functional requirements.

Functional specification (pseudocode)

```
get ASCII byte
sum bits 0 thru 6
put bit(0) of sum in bit(7) of ASCII byte
put ASCII byte somewhere
```

Now define how to sum bits 0 thru 6

```
set counter to 0           ;bit pointer
set sum to 0               ;sum of bits

loop:
sum=sum+byte[counter]     ;sum up bits 0...6
                           ;byte is ASCII character being
                           ;processed

counter=counter+1
if counter<7 goto loop
byte[7]=sum[bit0]         ;if sum[bit0] is 1 the sum is odd
                           ;if sum[bit1] is 0 the sum is even
                           ;this program generates even
                           ; parity
```

For even parity, if bits 0 thru 6 sum to an odd number then set bit #7 to 1 to make the parity even. If you wanted to change the program to odd parity, you simply need to change the last line of the pseudocode.

Examples:

If the sum of the character's bits is an odd number then the parity bit must be set to 1.

1	0	0	0	0	1	1	1
---	---	---	---	---	---	---	---

If the sum of the character's bits is an even number then the parity bit must be set to 0.

0	0	0	0	0	1	1	0
---	---	---	---	---	---	---	---

MC68000 assembly code for parity program:

```
main_loop EQU      *
* could have also used i/o to get data from keyboard
        MOVE.B     $1000,D1          ;get ASCII byte from $1000
* used quick instructions but not necessary
        MOVEQ      #0,D0             ;clear counter
        MOVEQ      #0,D2             ;clear sum

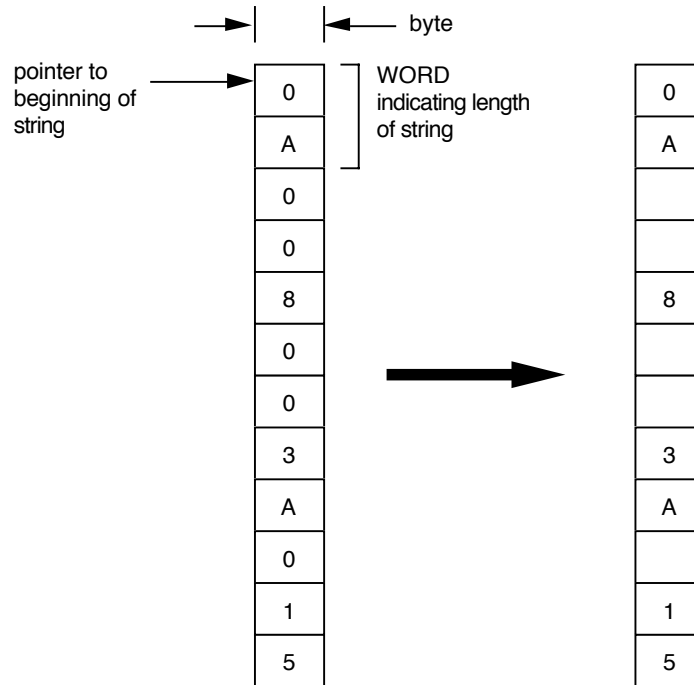
SUM      BTST.B     D0,D1             ;test D0-th bit of D1, sets Z-bit
        BEQ        SKIP_INCRE        ;if Z-bit=0 don't increment sum
        ADDQ       #1,D2             ;sum=sum+1

SKIP_INCRE
        ADDQ       #1,D0             ;increment counter

        MOVE       D0,D3             ;temp storage in D3
* subtract seven and compare to zero
        SUBQ       #7,D3             ;counter=7?
* could have used a compare instruction here
        BNE        SUM               ;No, sum more bits
        BCLR       #7,D1             ;Yes, clear parity bit
        BTST       #0,D2             ;get parity bit from sum[0]
        BEQ        PAR_SET           ;if parity bit=0 goto PAR_SET
        BSET       #7,D1             ;set parity bit to 1
PAR_SET  MOVE.B     D1,????          ;put ASCII byte somewhere
```

EXAMPLE: REPLACING 0's BY BLANKS PROGRAM

The correct way to design a program is by starting with your inputs, outputs and functional requirements.



Functional specification (pseudocode)

```

;define inputs
pointer=location of character string in memory
length=length of string (bytes) ;this will be contained in first
;word of string input
blank=' ' ;define a blank character
if (length=0) then quit ;if string length=0 do nothing

nextchar: ;basic loop for advancing to
;next character
if (char[pointer]≠'0') then ;if character is NOT a zero
    goto notzero ;then goto nonzero
char[pointer]=blank ;replace ASCII zero by blank
notzero:
length=length-1 ;decrement the char counter
if (length≥0) goto nextchar ;if more characters then repeat
    
```

What the program does is search for all the ASCII zeros in the string and replace them with blanks. This might be useful for eliminating leading zeros in a print routine.

SAMPLE PROGRAM

```

                ORG      $6000
START          DS.L      1           ;START is the address of the string
CHAR_0        EQU.B     '0'        ;define CHAR_0 as ASCII 0
BLANK         EQU.B     ' '        ;define BLANK as ASCII space

                ORG      $4000
begin         MOVEA.L   START,A0    ; set pointer to start of string,
                                ; cannot use LEA START

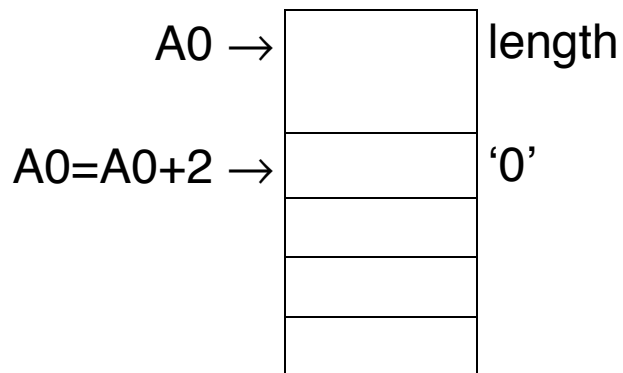
                MOVEQ    #BLANK,D1  ; put a blank in D1
                MOVE.W   (A0)+,D2   ; get length of string
                BEQ      DONE       ; if the string is of length zero
                                ; then goto DONE

NEXT_CHAR:
                MOVEQ    #CHAR_0,D0 ; put ASCII 0 into D0
                SUB.B    (A0)+,D0   ; compute '0'-current character
                BNE     NOT_ZERO    ; goto next char if non-zero
                MOVE.B   D1,-1(A0)  ; go back, get last byte and
                                ; replace it by ASCII zero

NOT_ZERO:
                SUBQ    #1,D2       ; decrement the character counter
                BPL     NEXT_CHAR   ; if count >=0 go to next character
                                ; otherwise quit

DONE          END        begin

```



EXAMPLE: LONG DIVISION USING REPEATED SUBTRACTION

Input, using HexIn, nonnegative numbers M and N where $N > 0$. Using repeated subtraction, find the quotient M/N and remainder.

Algorithm

Repeatedly subtract the divisor N from M ($M := M - N$). Count the number of iterations Q until $M < 0$. This is one too many iterations and the quotient is then $Q - 1$. The remainder is $M + N$, the previous value of M.

Pseudocode:

```
QUOTIENT:=0;
READLN(M);    {No error checking. Assume  $M \geq 0$ }
READLN(N);    {No error checking. Assume  $N \geq 0$ }
REPEAT
    QUOTIENT:=QUOTIENT+1;
    M:=M-N;
UNTIL M<0;
QUOTIENT:=QUOTIENT-1;
REMAINDER:=M+N;
```

Sample calculations:

Suppose $Q = \$0000$, $R = \$0000$

Start with $M = \$0015$, $N = \$0004$ {corresponds to $15/4 = 4$ w
/remainder=3}

Q=1: $M = M - N = \$0015 - \$0004 = \$0011$

Q=2: $M = M - N = \$0011 - \$0004 = \$000D$

Q=3: $M = M - N = \$000D - \$0004 = \$0009$

Q=4: $M = M - N = \$0009 - \$0004 = \$0005$

Q=5: $M = M - N = \$0005 - \$0004 = \$0001$

Q=6: $M = M - N = \$0001 - \$0004 = \$FFFD$

Since quotient is negative stop algorithm and back up one.

$Q = Q - 1 = 6 - 1 = 5$;correct quotient

$R = M + N = \$FFFD + \$0004 = \$0001$;correct remainder

SAMPLE PROGRAM

```
                INCLUDE   io.s                ;contains the i/o routines
                ORG      $6000
START          MOVE.W   #0,D2                ;quotient in D2, set to zero
GETM          JSR      HexIn                ;get M, put in D0
                TST.W   D0                  ;test for M≥0
                BMI     GETM                ;if M<0 get another M
                MOVE.W  D0,D1                ;put M in D1
GETN          JSR      HexIn                ;get N, put in D0
                TST.W   D0                  ;test for N>0
                BPL     LOOP                ;if N>0, start calculations
                BRA     GETN                ;if N≤0 get another N
LOOP          ADDI.W   #1,D2                ;increment the quotient
                SUB.W   D0,D1                ;compute M-N
                BPL     LOOP                ;branch back if M not negative,
                                           ;corresponds to doing another
                                           ;division
RESULT       SUBI.W   #1,D2                ;decrement the quotient
                ADD.W   D0,D1                ;set remainder
                MOVE.W  D2,D0                ;move quotient to D0
                JSR     HexOut               ;display quotient
                MOVE.W  D1,D0                ;move remainder to D0
                JSR     HexOut               ;display remainder
                JSR     NewLine              ;advance to next line
                TRAP    #0                  ;trick to end program
                END     START
```


EXAMPLE: Tests for Signed and UnSigned Overflow

Description:

Enter two 16-bit numbers and compute their sum. The addition operation sets the CCR bits. These bits are then read from the SR into the least significant word of D0 using the MOVE SR,Dn instruction. After isolating the C and V bits in D0, a message indicating if overflow has occurred is printed.

Pseudocode:

```
READLN(M);           /*No error checking. Assume M≥0*/
READLN(N);           /*No error checking. Assume N≥0*/
M:=M+N;
D0:=SR;              /*put the value of the SR into D0*/
D0:=D0&&0x0003;      /*Clear bits 2-15 by ANDing with $0003*/
WRITELN(D0);         /*Write out D0*/
SWITCH (D0) {
    CASE 1:  WRITELN('NO OVERFLOW'); BREAK;
    CASE 2:  WRITELN('ONLY UNSIGNED OVERFLOW');
              BREAK;
    CASE 3:  WRITELN('ONLY SIGNED OVERFLOW'); BREAK;
    CASE 4:  WRITELN('SIGNED AND UNSIGNED
              OVERFLOW'); BREAK;
    DEFAULT;
}
```

MASKing:

ANDI.W #D0 masks bits 0-1

$0003_{16} = 0000\ 0000\ 0000\ 0011_2$

$(D0) = \underline{\text{xxxx}\ \text{xxxx}\ \text{xxxx}\ \text{xxxx}_2}$

$(D0) = 0000\ 0000\ 0000\ 00xx_2$

Since the AND operates according to $0 \cdot x = 0$ and $1 \cdot x = x$ the result contains only whatever was in bits 0 and 1 — all other bits were set to

zero. Basically we masked out bits 0 and 1; hence the name, masking.

SAMPLE PROGRAM

```

                INCLUDE   io.s                ;contains the i/o routines
                ORG      $6000
START          JSR      HexIn                ;get M, put in D0
                MOVE.W   D0,D1              ;put M in D1
                JSR      HexIn                ;get N, put in D0
                ADD.W    D0,D1              ;D0:=M+N
                MOVE     SR,D0              ;get contents of SR
                ANDI.W   #$0003,D0         ;clears bits 2-15
                JSR      HexOut              ;display C and V bits
                LEA     OVRFLSTR,A1         ;base address of output messages
                ADD.W    D0,D0              ;compute 4*D0 by adding D0 to
                                           ;itself twice
                ADD.W    D0,D0              ;faster than a multiply
                ADDA.L   D0,A1              ;add message offset to base
                                           ;address
                MOVEA.L  (A1),A0           ;set (A1) to start address of
                                           ;message
                MOVE.W   #28,D0            ;each string has 28 characters
                                           ;(bytes)
                JSR      StrOut              ;string output routine
                JSR      NewLine            ;advance line
                TRAP     #0                 ;exit to debugger

OVRFLSTR      DC.L      NO_OVR,USGNOVR,SGNOVR,DUALOVR
NO_OVR        DC.B      'NO OVERFLOW      '
USGNOVR       DC.B      'ONLY UNSIGNED OVERFLOW '
SGNOVR        DC.B      'ONLY SIGNED OVERFLOW '
DUALOVR       DC.B      'UNSIGNED AND SIGNED OVERFLOW'
                END      START

```

HOW DOES PROGRAM IMPLEMENT SWITCH:

```
LEA    OVRFLSTR,A1
```

loads the base address of the table of messages

D0 can only have the values

<u>D0</u>	<u>V</u>	<u>C</u>
0	0	0
1	0	1
2	1	0
3	1	1

Multiply D0 by 4 to make these values in D0 correspond to the message since

```
OVRFLSTR    DC.L
```

```
NO_OVR,USGNOVR,SGNOVR,DUALOVR
```

places the beginning addresses of the messages in consecutive long words beginning at OVRFLSTR.

Use

```
MOVEA.L (A1),A0
```

to get the starting address of the correct message into A0

NOTE:

```
MOVEA.L A1,A0
```

will simply place the address of the address of the message into A0 which is NOT what was wanted.

The instruction

```
LEA    0(A1,D0.W),A0
```

would have also worked by directly adding the offset