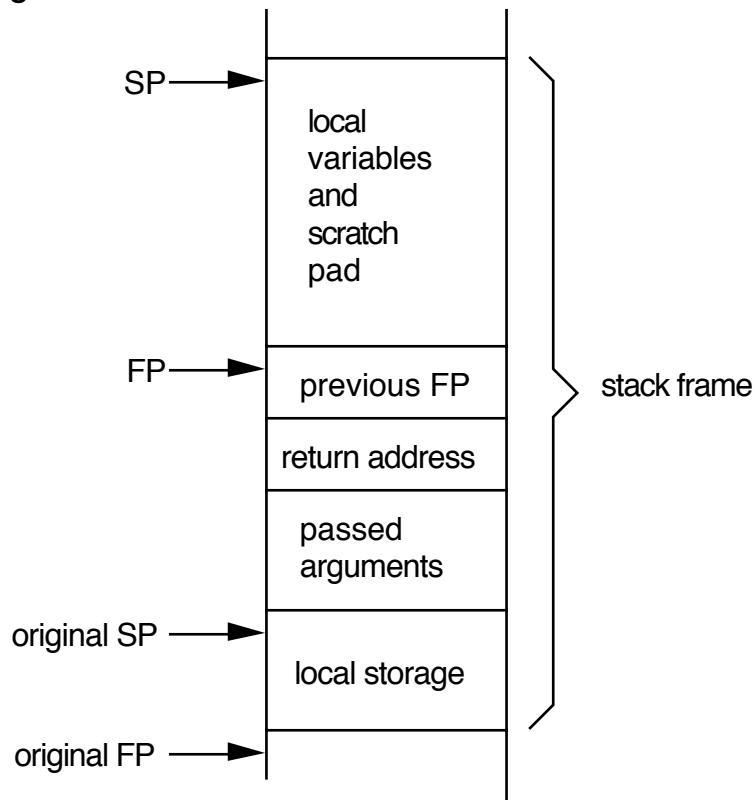


STACK FRAMES

The MC68000 provides two special instructions to allocate and deallocate a data structure called a frame in the stack to make subroutines easier to code.

general structure of a frame:



where register An is used as the argument pointer.

LINK An,d

1. put An at -(SP) Example:
decrement stack pointer and put
A0 on the stack.

2. put SP into An Example:
set A0 to point to this value.

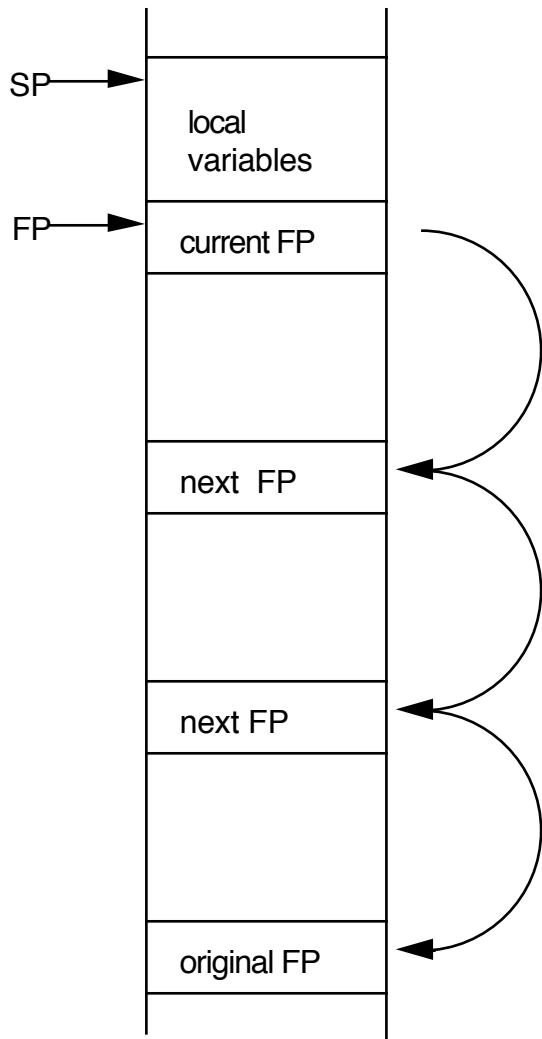
3. change SP-d to SP, i. e.
decrement the SP

UNLK An

1. An → SP, change the value of
the SP to that contained in An

2. $(SP)+ \rightarrow A_n$, put that value on the stack into A_n and deallocate that stack space.

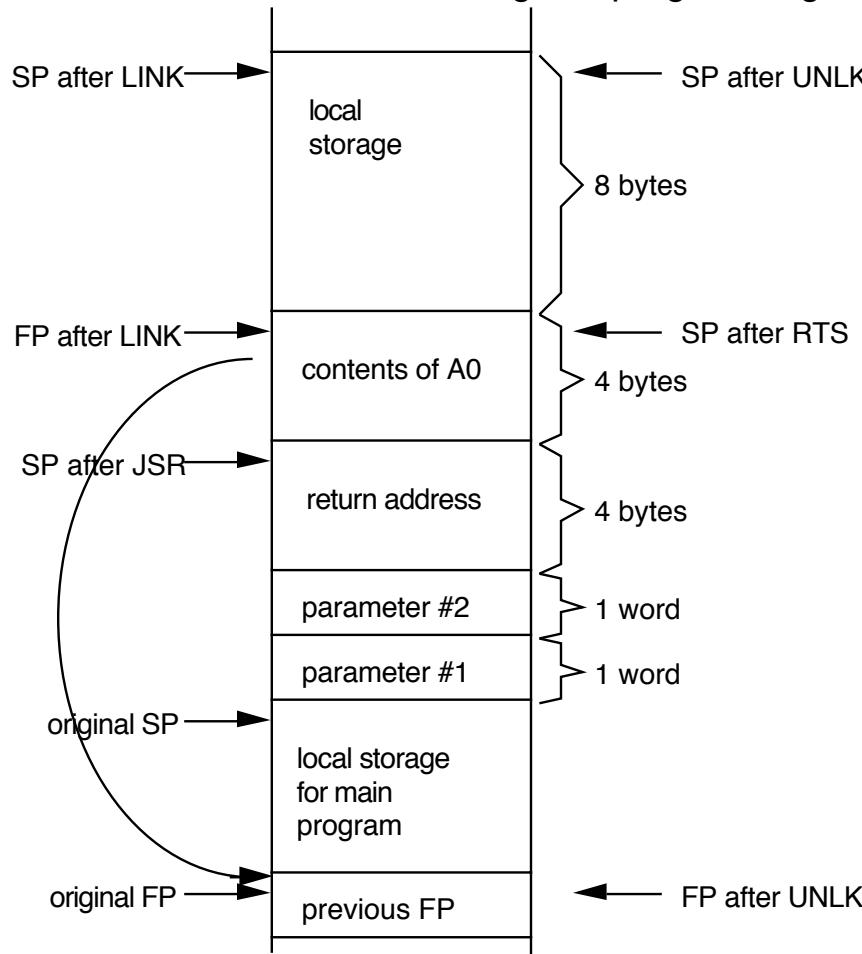
Return addresses and passed arguments are always positive relative to the frame pointer (FP).



Example:

	MOVE.W	D0,-(SP)	;push parameter #1 onto stack
	MOVE.W	D1,-(SP)	;push parameter #2 onto stack
	JSR	SBRT	;jump to subroutine SBRT
SBRT	LINK	A0,-#\$8	;establish FP and local storage
	.		
	.		
	.		
	MOVE.W	10(A0),D5	;retrieve parameter #1
	.		
	.		
	.		
	UNLK	A0	;FP for the calling routine re-established. Deallocate stack frame
	RTS		;return

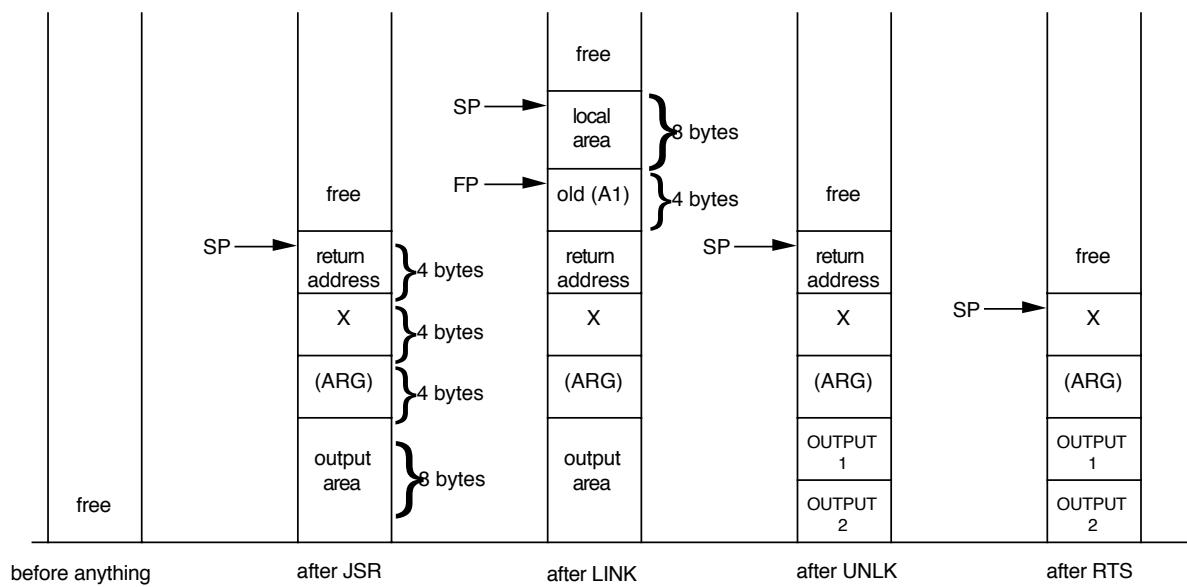
What the stack looks like during this program segment:



Note that the FP is stored in A0.

EXAMPLE:

ARG	DC.L		;number
N	EQU	8	;8 bytes for output
M	EQU	8	;8 bytes for local variables
	ADD.L	#-N,SP	;put output area on stack
	MOVE.L	ARG,-(SP)	;put argument on stack
	PEA	X	;put address of data table on stack
	JSR	SUBR	;goto subroutine
	ADDA	#8,SP	
	MOVE.L	(SP)+,D1	;read outputs
	MOVE.L	(SP)+,D2	
	.		
	.		
	.		
SUBR	LINK	A1,#-M	;save old SP
	.		
	.		
	.		
	MOVE.L	LOCAL1,-4(A1)	;save old variables
	MOVE.L	LOCAL2,-8(A1)	;
	.		
	.		
	.		
	ADD.L	#1,-4(A1)	;change a local variable
	MOVEA.L	8(A1),A2	;get X
	.		
	.		
	.		
	MOVE.L	OUTPUT,16(A1)	;push an output
	.		
	.		
	.		
	UNLK	A1	
	RTS		
LOCAL1	DC.L	\$98765432	;local variables
LOCAL2	DC.L	\$87654321	
OUTPUT	DC.L	'ADCB'	output value



Program to compute the power of a number using a subroutine.
 Power MUST be an integer. A and B are signed numbers.
 Parameter passing using LINK and UNLK storage space on the stack.

```

MAIN  LINK      A3,#-6      ;sets up SP
      MOVE      A,-2(A3)
      MOVE      B,-4(A3)
      JSR       POWR      ;call subroutine POWR
      LEA       C,A5
      MOVE      -6(A3),(A5)
      UNLK      A3

ARG   EQU      *
A     DC.W     4
B     DC.W     2
C     DS.W     1

POWR  EQU      *
      MOVE      -2(A3),D1      ;put A into D1
      MOVE      -4(A3),D2      ;put B into D2
      MOVE.L    #1,D3      ;put starting 1 into D3
LOOP   EQU      *
      SUBQ      #1,D2      ;decrement power
      BMI       EXIT      ;if D2-1<0 then quit NOTE: this
                           gives us A**0=1
      MULS      D1,D3      ;multiply out power
      BRA      LOOP      ;and repeat as necessary
EXIT   EQU      *
      MOVE      D2,-6(A3)    ;C=(D3)
      RTS

END    MAIN

```

	(D2)	2 bytes
	B	2 bytes
*FP→	A	2 bytes
SP→	value of A3	

*fixed while the SP changes

Better way.

```
MAIN  MOVEA.L  SP,A3
      MOVE      A,-(SP)
      MOVE      B,-(SP)
      ADD.L    #2,SP      ;save output area
      JSR       POWR      ;call subroutine POWR
      LEA       C,A5
      MOVE     -6(A3), (A5) ;put answer somewhere

ARG   EQU      *
A    DC.W     4
B    DC.W     2
C    DS.W     1

POWR  EQU      *
      LINK     A3,#-6
      MOVE     10(A3),D1    ;put A into D1
      MOVE     12(A3),D2    ;put B into D2
      .
      .
      .
      MOVE     D2,8(A3)   ;C=(D3)
      UNLK     A3
      RTS

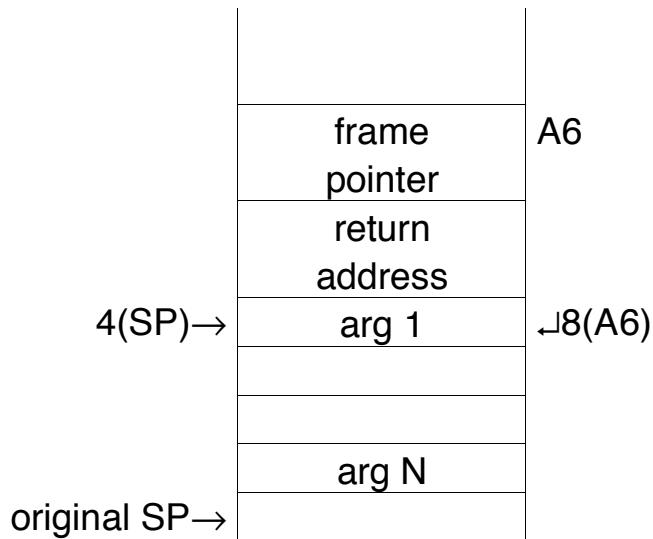
END   MAIN
```

SP→	C	save 6 bytes for next set of variables
	B	
	A	
FP (in A3)→	previous FP	4 bytes
	return address	4 bytes
	reserved for C	2 bytes
	B	2 bytes
	A	
original SP,FP→		

Calling conventions for C or Pascal

Arguments are pushed onto the stack in the reverse order of their appearance in the parameter list.

Just after a subroutine call:

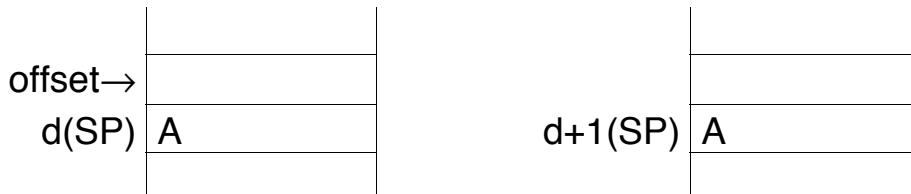


If the function begins with a

LINK A6,#

High level language always generates **LINK A6,#** instructions

All arguments occupying just a byte in C are converted to a word and put in the low byte of the word, i.e.



Result, if any, is returned in D0 for function calls.

IT IS THE PROGRAMMER'S RESPONSIBILITY TO REMOVE THE ARGUMENTS FROM THE STACK.

The C calling sequence looks like this:

```
MOVE    ___,-(SP) ;last argument
•
•
•
MOVE    ___,-(SP) ;first argument
JSR     FUNCT
ADD    #N,SP      ;total size of arguments
```

Subroutine functions:

```
LINK   A6,#N
•
•
•
MOVE   ...,D0
UNLK   A6
RTS
```

The Pascal calling sequence pushes arguments in left to right order, then calls the function. The result if any is left on the stack. An example looks like this:

```
SUB    #N,SP      ;save space for result
MOVE   ...,-(SP)   ;push first argument onto stack
•
•
•
MOVE   ...,-(SP)   ;last argument
JSR    FUNCT
MOVE   (SP)+,...  ;store result
```

Subroutine code:

```
LINK      A6,#N
.
<code>
.
UNLK      A6
MOVE      (SP)+,A0 ;return address
ADD       #N,SP    ;total size of arguments
MOVE      ..., (SP) ;store return result
JMP       (A0)
```

Symbols defined in assembly routines with the DS directive and exported using XDEF and XREF can be accessed from C as external variables. Conversely, C global variables can be imported and accessed from assembly using the XREF directive.

Miscellaneous comments about subroutines.

Parameter passing via MOVEM (move multiple registers)

If you have a small assembly language program this instruction allows you to save the values of registers NOT used to pass parameters.

Example:

```
SUBRTN EQU      *
          MOVEM    D0-D7/A0-A6,SAVBLOCK
          .
          .
          .
          MOVEM    SAVBLOCK,D0-D7/A0-A6
```

where SAVBLOCK is local memory. This is bad practice since SAVBLOCK can be overwritten by your program.

MOVEM has two forms

```
MOVEM    register_list,<ea>
MOVEM    <ea>,register_list
```

More common to save registers on stack

```
SUBRTN EQU      *
          MOVEM    D0-D7/A0-A6,-(SP)
          .
          .
          .
          MOVEM    (SP)+,D0-D7/A0-A6
          RTS
```

MOVEM is often used for re-entrant (subroutines that can be interrupted and re-entered) procedures.

The MOVEM instruction always transfers contents to and from memory in a predetermined sequence, regardless of the order used to specify them in the instruction.

address register indirect with pre- transferred in the order A7→A0,
decrement then D7→D0

for all control modes and address register indirect with post-increment transferred in reverse order
D0→D7, then A0→A7

This allows you to easily build stacks and lists.

Six methods of passing parameters:

1. Put arguments in D0 thru D7 before JSR (good only for a few arguments)
2. Move the addresses of the arguments to A0-A6 before JSR
3. Put the arguments immediately after the call. The argument addresses can be computed from the return address on the stack.
4. Put the addresses of the arguments immediately after the call in the code.
5. The arguments are listed in an array. Pass the base address of the array to the subroutine via A0-A6.
6. Use LINK and UNLK instructions to create and destroy temporary storage on the stack.

JUMP TABLES

- are similar to CASE statements in Pascal
- used where the control path is dependent on the state of a specific condition

EXAMPLE:

This subroutine calls one of five user subroutines based upon a user id code in the low byte of data register D0. The subroutine effects the A0 and D0 registers.

	RORG	\$1000	;causes relative addressing (NOTE 1)
SELUSR	EXT.W	D0	;extend user id code to word
	CHK	#4,D0	;invalid id code ? (NOTE 2)
	LSL	#2,D0	;NO! Calculate index=id*4 since all long word addresses
	LEA	UADDR,A0	;load table addresses
	MOVEA.L	0(A0,D0.W),A0	;compute address of user specified subroutine and put correct calling address into A0
	JMP	(A0)	;jump to specified routine
	•		
	•		
	•		
UADDR	DC.L	USER0,USER1,USER2,USER3,USER4	

NOTES:

1. The RORG is often used when mixing assembly language programs with high level programs. It causes subsequent addresses to be relative.
2. The CHK is a new instruction. In this case it checks if the least significant word of D0 is between 0 and 4 (2's complement). If the word is outside these limits, an exception through vector address \$10 is initiated. The CHK instruction checks for addresses outside assigned limits and is often used to implement subscript checking.

EXAMPLE RECURSIVE PROCEDURE USING STACK

```
DATA EQU      $6000
PROGRM EQU     $4000

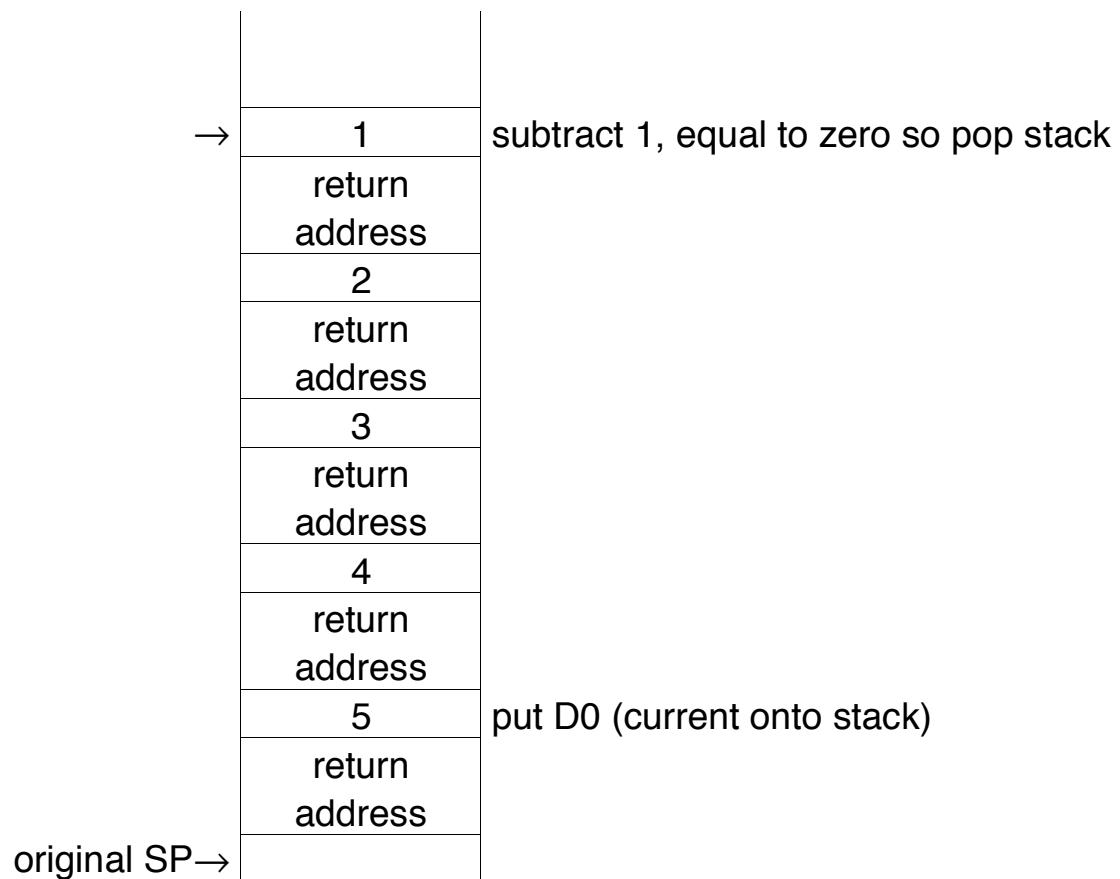
                ORG      DATA
NUMB    DS.W     1          ;number to be factorialized
F_NUMB DS.W     1          ;factorial of input number

                ORG      PROGRM
MAIN    MOVE.W   NUMB,D0    ;get input number
        JSR       FACTOR    ;compute factorial
        MOVE.W   D0,F_NUMB  ;save the answer
```

* SUBROUTINE FACTOR

- * PURPOSE: Determine the factorial of a given number.
- * INPUT: D0.W = number whose factorial is to be computed
- * $0 \leq D0.W \leq 9$
- * OUTPUT: D0.W = factorial of input number
- * REGISTER USAGE: No registers except D0 effected
- * SAMPLE CASE: INPUT: D0.W=5
- * OUTPUT: D0.W=120

```
FACTOR  MOVE.W   D0,-(SP)   ;push current number onto
                    stack
        SUBQ.W   #1,D0      ;decrement number
        BNE.S    F_CONT    ;not end of factorial
                    computations
        MOVE.W   (SP)+,D0    ;factorial=1
        BRA.S    RETURN
F_CONT   JSR       FACTOR
        MULU     (SP)+,D0
RETURN  RTS
```



EXAMPLE

This is a simplified version of TUTOR's "DF" command. It uses the stack to display register contents.

```
START    MOVEM.L TESTREGS,D0-D7/A0-A6 ;assign values to
          registers
        MOVE,L  #-1,-(SP)      ;put something on stack
        JSR     PRINTR       ;print all registers
        MOVE.L  (SP)+,D0      ;retrieve it
        ADDQ.L  #1,D0         ;null it
        JSR     PRINTR       ;print them all again
        TRAP   #0             ;stop program

SAVESP   EQU    60

PRINTR               ;data for PRINTREGS
RMSG$:
        DC.B   ' D0  D1  D2  D3  D4  D5',0
        DC.B   ' D6  D7  A0  A1  A2  A3',0
        DC.B   ' A4  A5  A6  SP  SR  PC',0

;
;----- 55 characters long -----|
SPACES   DC.B   ' ',0           ;2 blanks
CONBUF   DS.B   10
ENDLINE  DC.B   $0D,$0A,0
; data for program
CH       DS.B   1
          DS.W   1
TSTREG   DC.L   1,2,3,4,5,6,7,8,$A,$AA,$AAA
          DC.L   $AAAA,$AAAAAA,$AAAAAAA,$AAAAAAA
          END
```

	PRINTR	MOVE.W SR,-(SP)	;save SR on stack
		PEA 6(SP)	;save original SP on stack
		MOVEM.L D0-D7/A0-A6,-(SP)	;save all regular registers
		MOVEQ #2,D4	;D1 counts # of rows in printout
		MOVEA.L SP,A1	;use A1 to point to beginning of data
		LEA RMSGS,A2	;use A2 to point to row headings
MLOOP			;output routine for heading
		MOVEA.L A2,A0	;set pointer to beginning of header to be printed
		JSR PrintString	;output heading
		MOVEQ #5,D5	;output six registers this line
RLOOP		TST.W D4	;tests for SR to be printed
		BNE.S NOT_SR	;SR requires special routine
		CMP.W #1,D5	;as it is only word length
		BNE.S NOT_SR	;register
		LEA SPACES,A0	;load addresses of spaces
		JSR PrintString	;print spaces with no new line
		MOVE.W (A1)+,D0	;put SR word into D0
		JSR PNT4HX	;unimplemented routine to convert 4 hex digits in D0 to ascii code for printing
		JSR PrintString	;print hex contents
		LEA SPACES,A0	;load address of spaces
		JSR PrintString	;print them with no line feed
		BRA.S ENDRPL	
NOT_SR		MOVE.L (A1)+,D0	;put register contents into D0
		JSR PNT8HX	;unimplemented routine to convert 8 hex digits in D0 to ascii code for printing

ENDRPL	DBF	D5,PRLOOP	;decrement register counter, started at 5
	LEA	ENDLINE,A0	;print CR+LF
	JSR	PrintString	
	ADDA.L	#55,A2	;increment heading pointer
	DBF	D4,MLOOP	;goto another line
	MOVEM.L	(SP)+,D0-D7/A0-A6	
	ADDQ.W	#4,SP	;skip over A7 to point to SR
	RTR		;return and restore registers

